

Colloquium Tonight! 5pm on Zoom

Closure properties for regular languages

If L is a regular language, then
 $\bar{L}$ is also a regular language.

regular langs are "closed under complement"

If we have a DFA for L ,

the DFA for $\bar{L}$ is obtained by swapping
all accepting/rejecting states.

Intersections

Thm If L_1 & L_2 are regular languages, then
 $L_1 \cap L_2$ is a regular language.

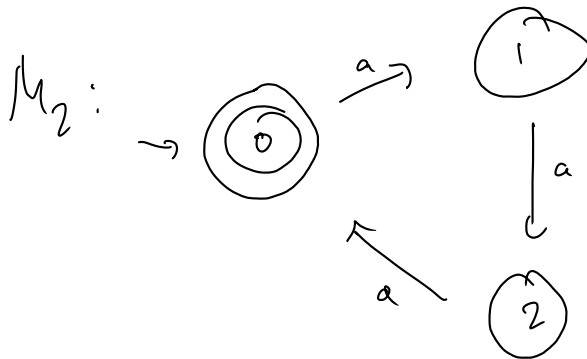
On DFAs, this is done using "the product construction"

Ex $\Sigma = \{a\}$

$$L_1 = \{a^n \mid n \text{ is even}\}$$

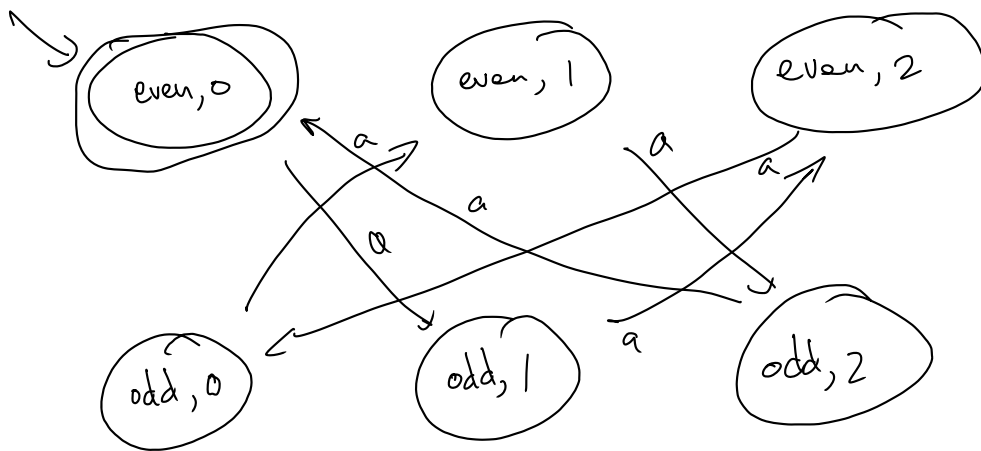
want $L_1 \cap L_2$

$$L_2 = \{a^n \mid n \text{ is div. by } 3\}$$



We want to combine these two "with AND"

Combine the states in pairs:

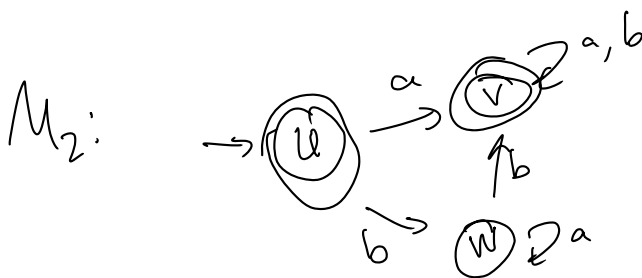
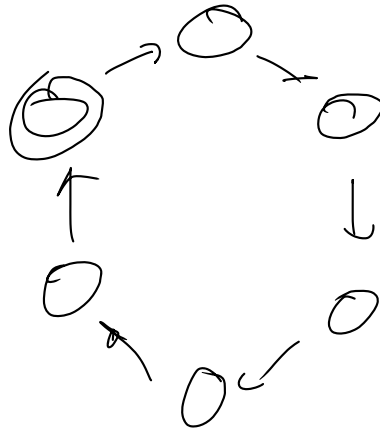


The states are pairs of states from the original machines,
 transitions apply the letter to each part of pair,
 starting state is the pair obtained from the
 starting states of M_1 & M_2

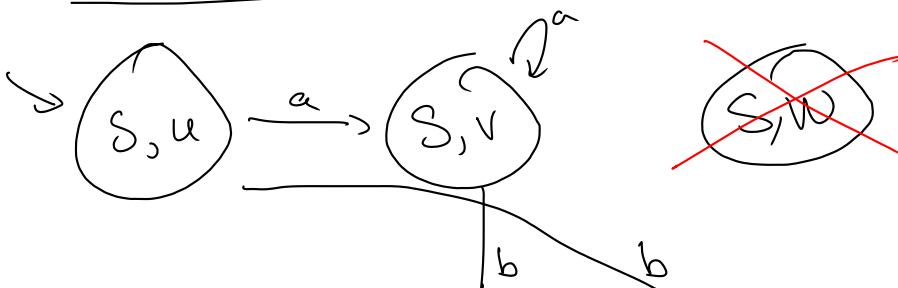
Accepting states are any pair of states

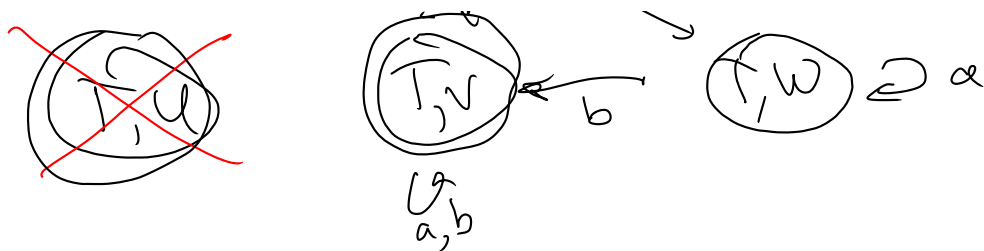
which are both accepting in M_1 & M_2 .

ext is equivalent to

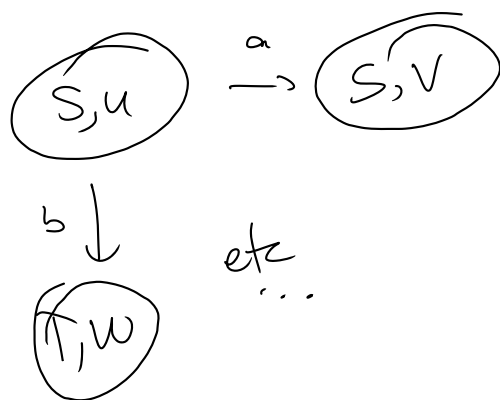


Make a DFA
for the intersection.





Some pairs may be unreachable from the starting state.



"Prove the product construction"

Say $M_1 = (Q, \Sigma, \delta_1, q_0, F_1)$

$M_2 = (R, \Sigma, \delta_2, r_0, F_2)$

Then the product construction results in:

$$M_3 = (Q \times R, \Sigma, \delta, (q_0, r_0), F_1 \times F_2)$$

$$\text{where } \delta((q, r), a) = (\delta_1(q, a), \delta_2(r, a))$$

Our proof should show that

$$L(M_3) = L(M_1) \cap L(M_2)$$

Remember

$$L(M) = \{x \in \Sigma^* \mid \delta^*(q_0, x) \in F\}$$

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$$L(M_3) = \{x \in \Sigma^* \mid \delta^*((q_0, r_0), x) \in F_1 \times F_2\}$$

$$= \{x \in \Sigma^* \mid (\delta_1^*(q_0, x), \delta_2^*(r_0, x)) \in F_1 \times F_2\}$$

$$= \{x \in \Sigma^* \mid \delta_1^*(q_0, x) \in F_1 \text{ and } \delta_2^*(r_0, x) \in F_2\}$$

$$= \{x \in \Sigma^* \mid \delta_1^*(q_0, x) \in F_1\} \cap \{x \in \Sigma^* \mid \delta_2^*(r_0, x) \in F_2\}$$

$$= L(M_1) \cap L(M_2)$$

Shown!