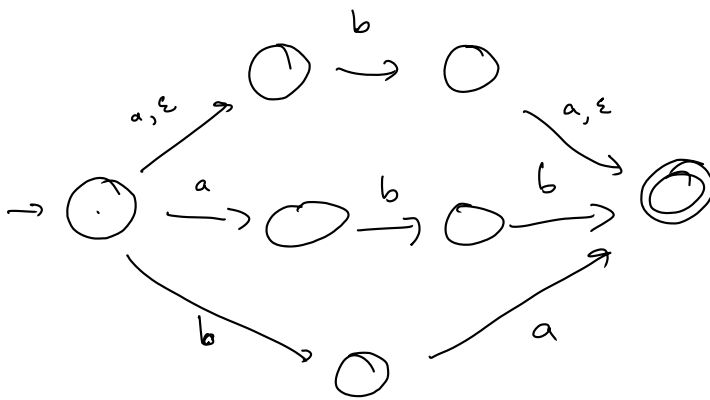


Last time: $\text{regex} \rightarrow \text{NFA}$

Can we do $\text{NFA} \rightarrow \text{regex}$?

Yes!

Make this into regex:



What are the paths from start to accepting?

write each path as regex, combine with +.

bottom: ba

middle: abb

top:

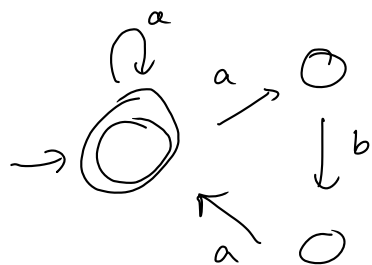
$aba + ab + ba + b$

$(a + \epsilon)b(a + \epsilon)$

The NFA is equivalent to

$(a + \epsilon)b(a + \epsilon) + abb + ba$

When there are loops, we must use *



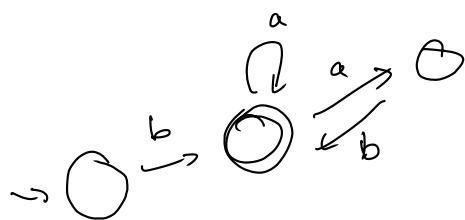
which loops can we do?

the a loop would be a^*

the other would be $(aba)^*$

$$(a + aba)^*$$

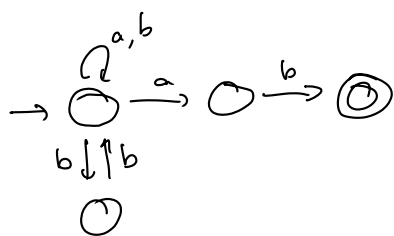
"sum the loops inside a *"



$$b(a + ab)^*$$



$$b(a + aa^*b)^*$$



$$((a+b) + bb)^* ab$$

$$(a + b + bb)^* ab$$

$$(a+b)^* ab$$

Any NFA can be converted to a regex.

We can go back & forth
 $\text{regex} \leftrightarrow \text{NFA}$

Then Kleene's Theorem

L is a regular language iff

L can be generated by regex.

There are non-regular languages,

$\{a^n b^n \mid n \in \mathbb{N}\}$ is not regular

A general method to show a language
is non-regular.

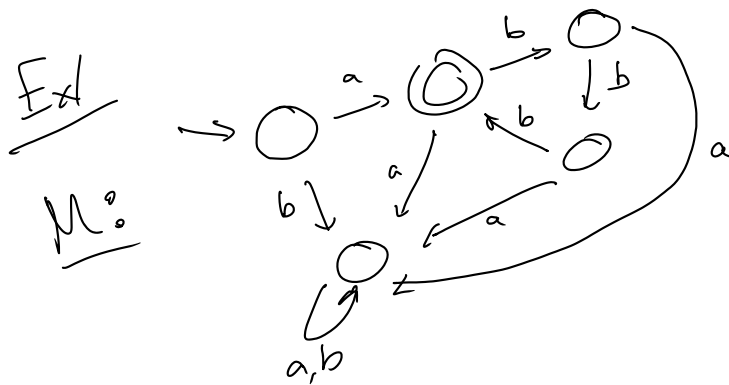
Involves the derivative!

The Brzozowski Derivative
(not in our book)

book uses The Pumping Lemma

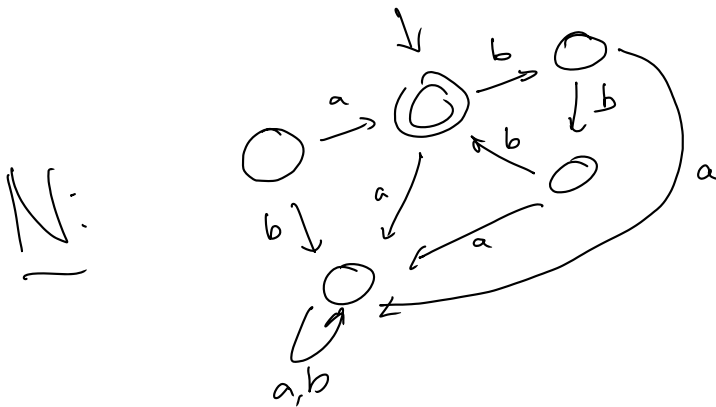
The derivative is related this question:

What happens to the language of
a NFA if I move the starting
state?



Here, $L(M) = \{ a(b^3)^n \mid n \in \mathbb{N} \}$

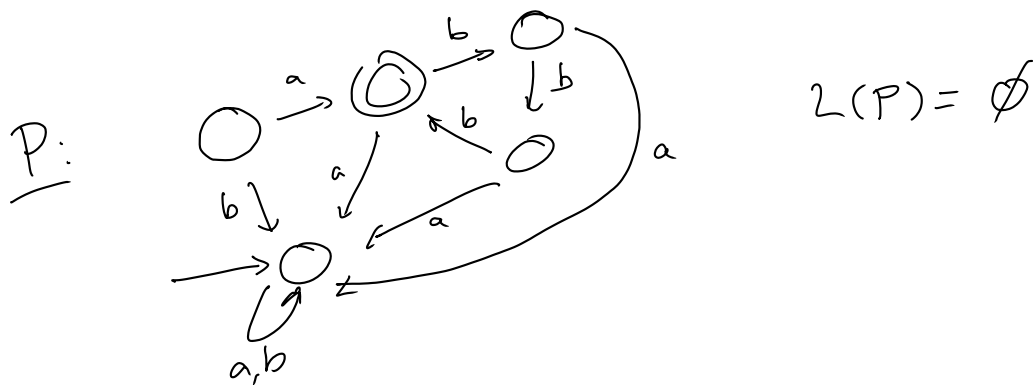
What if we move the start state thru the a edge?



$$L(N) = \{ (b^3)^n \mid n \in \mathbb{N} \}$$

Moving the start state thru the a edge
deletes a from the start of all our
strings.

Move it the other way:



We are skipping over any initial b's.

but $L(M) = \{ a (b^3)^n \}$ has no
initial b's, so we end up
with nothing.

What we're doing by "skipping initial a's" is
really:

$$\{ x \mid ax \in L \}$$

This is the B. derivative of L with
respect to a .

Def $\frac{d}{da} L = \{x \mid ax \in L\}$

For $L = \{a(b^3)^n\}$,

$$\frac{d}{da} L = \{(b^3)^n\}$$

$$\frac{d}{db} L = \emptyset$$

If L was the set of all English words,

what does $\frac{d}{da} L$ look like?

it contains strings like: how since allow $\in L$
 pe
 advertisement
 e