

last time: CFG $\rightarrow$ stack:

$$S \rightarrow aSb \mid T$$

$$T \rightarrow bT \mid \epsilon$$



read	pop	push
ϵ	S	aSb
ϵ	S	T
ϵ	T	bT
ϵ	T	ϵ
a	a	ϵ
b	b	ϵ

We can convert stack $\rightarrow$ grammar

A stack rule:

read	pop	push
x	y	z

becomes: $y \rightarrow xz$

$$\{x \mid \#a's = \#b's\}$$

	read	pop	push
1	a	S	XS
2	a	X	XX
3	a	Y	ϵ
4	b	S	YS
5	b	Y	YY
6	b	X	ϵ
7	ϵ	S	ϵ

Grammar:

1. $S \rightarrow aXS$
2. $X \rightarrow aXX$
3. $Y \rightarrow a$
4. $S \rightarrow bYS$
5. $T \rightarrow bYY$
6. $X \rightarrow b$
7. $S \rightarrow \epsilon$

$$S \rightarrow aXS \mid bYS \mid \epsilon$$

$$X \rightarrow aXX \mid b$$

$$Y \rightarrow bYY \mid a$$

Derive abba

stack: $(\underline{abba}, S) \mapsto_1 (\underline{bba}, XS) \mapsto_2 (\underline{ba}, S) \mapsto_4 (\underline{a}, YS)$
 $\mapsto_3 (\underline{\epsilon}, S) \mapsto_7 (\underline{z}, z)$

grammar: $\underline{S} \Rightarrow_1 \underline{aXS} \Rightarrow_2 \underline{abS} \Rightarrow_4 \underline{abYS} \Rightarrow_3 \underline{abbaS} \Rightarrow \underline{abba}$

Our grammar steps always look like

(string ^{read} _{so far}) (stack ^{built} _{_____})
 $\uparrow$
 action happens here.

What action exactly?

if the stack rule is

	read	pop	push
	x	y	z,

then $(\underline{\text{string}}) \underline{y}(\underline{\text{stack}})$ will become:

$(\text{string } x) z(\text{stack})$

this was a grammar substitution $y \rightarrow xz$

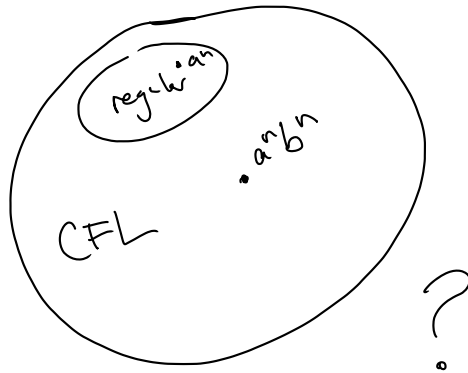
Last time: CFG $\rightarrow$ stack.

This time: stack $\rightarrow$ CFG

So CFG & stack machines have the same capabilities.

Then L is context-free iff

L is the language of some stack machine.



Are there any
simple non-CF
languages?

Classic example: $a^n b^n c^n$ is not a CFL

Failed attempt: $a^n b^n$ is $S \rightarrow aSb \mid \epsilon$

for $a^n b^n c^n$: ~~$S \rightarrow aSbSc \mid \epsilon$~~ makes $a^n (bc)^n$

~~$S \rightarrow aSbSc \mid \epsilon$~~

~~$S \Rightarrow aSbSc \Rightarrow aaSbScbaSbSc$
 $\Rightarrow a^2 bcbabc$~~

$S \rightarrow aSTR \mid \epsilon$
 $T \rightarrow b$

$$P \rightarrow C$$

Later (maybe never)

we'll show why $a^n b^n c^n$ is not CFL.

A grammar for $a^n b^n c^n$:

$$S \rightarrow aSBc \mid abc \mid \varepsilon$$

$$cB \rightarrow Bc \quad \leftarrow \text{B can commute with c}$$

$$bB \rightarrow bb \quad \leftarrow \text{B can become b, if it's next to b.}$$

$$S \Rightarrow aSBc \Rightarrow aabc \underline{B}c$$

$$\Rightarrow aabBcc \Rightarrow aabbcc$$

Most real-world programming langs are not CFL.

The Python interpreter must analyze & balance
more than 2 things at once to
parse indents

Python is indented like:

XXX :

T XXX

T XXX

XX

XXX :

T XX :

TT XXXX

TT XXXX

or

TT XX

TTT XXX

TTT XXX

TTT XXX

↑
we have something like

$T^2(\sim)T^2(\sim)T^2(\sim)T^2(\sim)$

→
can't do this in a CFL.

def [name]([vars]): INDENT BLOCK OUTDENT

Turing Machines

In early 1900s math was being rewritten using axioms:

1880s - axioms in algebra

1890s - Hilbert geometry axioms

1900s - axioms for set theory

Axioms for everything?

Can all facts be proved based on axioms?

Is there one single master procedure to prove stuff from axioms?

Turing has the idea to turn all of math into some kind of mechanized procedure.